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**TERMINAL POINTS, BANDWIDTH AVAILABILITY AND BANDWIDTH UTILIZATION IN A  
VIRTUAL PRIVATE NETWORK (VPN)**

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## ABSTRACT

Efficiency, Reliability and Cost are the foremost three major criteria usually considered in a telecommunication network. In order to make better use of available network resources, there is a need for planning the bandwidth allocation to communication demands. In this study, a bandwidth allocation problem with three major constraints were simulated and solved for 50, 100, 150, 200 and 250 nodes. The study discovered that increasing bandwidth is not absolutely the only way to take care of the load on networks as the bandwidth availability may not be commensurate with bandwidth utilization.

**Key Words:** Bandwidth, Reliability, Efficiency, Cost, Bandwidth Availability, Bandwidth Utilization

## INTRODUCTION

Communication networks are expected to offer a wide range of services to an increasingly large number of users, with a diverse range of quality of service. This calls for efficient control and management of these networks. The key resource to manage in networks is bandwidth. Therefore, in order to make better use of available network resources, there is a need for planning the bandwidth allocation to communication demands, in order to set up routing tables (or any other route selection criterion) more purposefully. This can be achieved by the use of global information, including not only the available link capacities but also the expected traffic profile. This traffic profile may be given, as when setting up virtual private networks in an ATM backbone of a

provider, or estimated by objective traffic measurements [1, 2]. The goal of a VPN is to provide the organization with the same secure capabilities but at a much lower cost. Without proactive management, network capacity fills with inappropriate traffic and viruses, and the connection becomes ineffective. [2, 3, 4]

Several studies make use of Genetic Algorithm (GA) based techniques to solve network problems. The motivation behind GA's in nonlinear optimization problems is that the problem can be expressed such that natural evolution, as reported, can provide an attractive paradigm for implementing general nonlinear searches [5, 6, 7].

## METHODOLOGY

A dynamic mathematical model was developed and a Genetic Algorithm method was adopted for the optimization solution.

Terms	Notation
Minimum Required Bandwidth at node $n_i$ .	$b_{mi}$
Available Bandwidth at node $n_i$	$b_i$
Throughput Request at node $n_i$ .	$t_{ri}$
Allocated Throughput at node $n_i$ .	$t_{ai}$
Total Available Bandwidth for all the nodes at a given time.	$BW$
time (evaluation period)	$t$

$$\max \sum_{i=1}^k t_{ai} \quad (1)$$

$$\text{s.t.} \quad \sum_{i=1}^k b_i \leq BW \quad (2)$$

$$t_{ai} \geq t_{ri} \quad \forall i = 1, 2, 3, \dots, k \quad (3)$$

$$b_i \geq b_{mi} \quad \forall i = 1, 2, 3, \dots, k \quad (4)$$

The GA algorithm was implemented with Java programming language. The model was simulated for ten thousand (10,000) generations. Fifty (50), one hundred (100), one hundred and fifty (150), two hundred (200), and two hundred and fifty (250) terminal points/nodes were studied. Each of the nodes were simulated differently under varying bandwidths values of 64Kbps, 128Kbps, 256Kbps, 512Kbps, 1Mbps, 2Mbps, 4Mbps, and 8Mbps for a period of twenty four (24) hours.

## RESULTS

The bandwidth utilization tables and curves are as in tables 1 – 5 and figures 1 – 5 respectively.

Table 1: Daily Bandwidth Utilization for 50 Nodes

<b>Period</b>	<b>@ 64Kbps</b>	<b>@128Kbps</b>	<b>@256Kbps</b>	<b>@512Kbps</b>	<b>@1Mbps</b>	<b>@2Mbps</b>	<b>@4Mbps</b>	<b>@8Mbps</b>
<b>12 midnight</b>	<b>0.983276</b>	<b>0.994263</b>	<b>0.999985</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.784922</b>	<b>0.380473</b>	<b>0.190487</b>
<b>1a.m</b>	<b>0.99884</b>	<b>0.994263</b>	<b>0.999985</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.805693</b>	<b>0.395408</b>	<b>0.200751</b>
<b>2a.m</b>	<b>0.99884</b>	<b>0.999298</b>	<b>0.999985</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.822029</b>	<b>0.405922</b>	<b>0.207489</b>
<b>3a.m</b>	<b>0.99884</b>	<b>0.999298</b>	<b>0.999985</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.829897</b>	<b>0.410142</b>	<b>0.211594</b>
<b>4a.m</b>	<b>0.99884</b>	<b>0.999298</b>	<b>0.999985</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.831614</b>	<b>0.412102</b>	<b>0.211601</b>
<b>5a.m</b>	<b>0.99884</b>	<b>0.999298</b>	<b>0.999985</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.831614</b>	<b>0.412431</b>	<b>0.212073</b>
<b>6a.m</b>	<b>0.99884</b>	<b>0.999298</b>	<b>0.999985</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.834846</b>	<b>0.41379</b>	<b>0.213897</b>
<b>7a.m</b>	<b>0.99884</b>	<b>0.999298</b>	<b>0.999985</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.84403</b>	<b>0.415778</b>	<b>0.216415</b>
<b>8a.m</b>	<b>0.99884</b>	<b>0.999298</b>	<b>0.999985</b>	<b>0.99987</b>	<b>0.999985</b>	<b>0.848207</b>	<b>0.416536</b>	<b>0.216501</b>
<b>9a.m</b>	<b>0.99884</b>	<b>0.999298</b>	<b>0.999985</b>	<b>0.99987</b>	<b>0.999985</b>	<b>0.853672</b>	<b>0.418024</b>	<b>0.21683</b>
<b>10a.m</b>	<b>0.99884</b>	<b>0.999298</b>	<b>0.999985</b>	<b>0.99987</b>	<b>0.999985</b>	<b>0.862913</b>	<b>0.423918</b>	<b>0.21683</b>
<b>11a.m</b>	<b>0.999756</b>	<b>0.999298</b>	<b>0.999985</b>	<b>0.99987</b>	<b>0.999985</b>	<b>0.864573</b>	<b>0.424318</b>	<b>0.217609</b>
<b>12noon</b>	<b>0.999756</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.99987</b>	<b>0.999985</b>	<b>0.866718</b>	<b>0.424461</b>	<b>0.21806</b>
<b>1p.m</b>	<b>0.999756</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.99987</b>	<b>0.999985</b>	<b>0.86709</b>	<b>0.425706</b>	<b>0.21821</b>
<b>2p.m</b>	<b>0.999756</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.99987</b>	<b>0.999985</b>	<b>0.868235</b>	<b>0.427423</b>	<b>0.218225</b>
<b>3p.m</b>	<b>0.999756</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.99987</b>	<b>0.999985</b>	<b>0.868635</b>	<b>0.427423</b>	<b>0.218575</b>
<b>4p.m</b>	<b>0.999756</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.99987</b>	<b>0.999985</b>	<b>0.869408</b>	<b>0.428367</b>	<b>0.220048</b>
<b>5p.m</b>	<b>0.999756</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.99987</b>	<b>0.999985</b>	<b>0.869408</b>	<b>0.428452</b>	<b>0.220084</b>
<b>6p.m</b>	<b>0.999756</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.99987</b>	<b>0.999985</b>	<b>0.871496</b>	<b>0.428495</b>	<b>0.220184</b>
<b>7p.m</b>	<b>0.999756</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.99987</b>	<b>0.999985</b>	<b>0.871496</b>	<b>0.429039</b>	<b>0.220206</b>
<b>8p.m</b>	<b>0.999756</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.99987</b>	<b>0.999985</b>	<b>0.871496</b>	<b>0.42964</b>	<b>0.220807</b>
<b>9p.m</b>	<b>0.999756</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.99987</b>	<b>0.999985</b>	<b>0.871496</b>	<b>0.429668</b>	<b>0.220985</b>
<b>10p.m</b>	<b>0.999756</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.99987</b>	<b>0.999985</b>	<b>0.871496</b>	<b>0.429683</b>	<b>0.221</b>
<b>11p.m</b>	<b>0.999756</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.99987</b>	<b>0.999985</b>	<b>0.874157</b>	<b>0.432401</b>	<b>0.221343</b>

**Table 2: Daily Bandwidth Utilization for 100 Nodes**

<b>Period</b>	<b>@ 64Kbps</b>	<b>@128Kbps</b>	<b>@256Kbps</b>	<b>@512Kbps</b>	<b>@1Mbps</b>	<b>@2Mbps</b>	<b>@4Mbps</b>	<b>@8Mbps</b>
<b>12 midnight</b>	0.98053	0.998383	0.996094	0.998955	0.999355	0.999985	0.777469	0.400851
<b>1a.m</b>	0.98053	0.998383	0.996094	0.998955	0.999355	0.999985	0.791616	0.404263
<b>2a.m</b>	0.993347	0.999756	0.996094	0.99987	0.999813	0.999985	0.808511	0.406952
<b>3a.m</b>	0.993347	0.999756	0.999069	0.99987	0.99987	0.999985	0.815034	0.408947
<b>4a.m</b>	0.994263	0.999756	0.999069	0.99987	0.999928	0.999985	0.817051	0.410156
<b>5a.m</b>	0.994263	0.999756	0.999069	0.999985	0.999928	0.999985	0.8214	0.413668
<b>6a.m</b>	0.994263	0.999756	0.999069	0.999985	0.999985	0.999985	0.822487	0.414555
<b>7a.m</b>	0.994263	0.999756	0.999069	0.999985	0.999985	0.999985	0.828166	0.414591
<b>8a.m</b>	0.994263	0.999756	0.999069	0.999985	0.999985	0.999985	0.828881	0.418267
<b>9a.m</b>	0.994263	0.999756	0.999069	0.999985	0.999985	0.999985	0.829782	0.41954
<b>10a.m</b>	0.994263	0.999756	0.999756	0.999985	0.999985	0.999985	0.834618	0.419598
<b>11a.m</b>	0.994263	0.999756	0.999756	0.999985	0.999985	0.999985	0.83612	0.420105
<b>12noon</b>	0.994263	0.999756	0.999756	0.999985	0.999985	0.999985	0.837007	0.420156
<b>1p.m</b>	0.994263	0.999756	0.999756	0.999985	0.999985	0.999985	0.837994	0.422258
<b>2p.m</b>	0.994263	0.999756	0.999756	0.999985	0.999985	0.999985	0.841298	0.426385
<b>3p.m</b>	0.999725	0.999756	0.999756	0.999985	0.999985	0.999985	0.841498	0.428674
<b>4p.m</b>	0.999725	0.999756	0.999756	0.999985	0.999985	0.999985	0.84167	0.429132
<b>5p.m</b>	0.999725	0.999756	0.999756	0.999985	0.999985	0.999985	0.842428	0.429139
<b>6p.m</b>	0.999725	0.999756	0.999756	0.999985	0.999985	0.999985	0.849295	0.429518
<b>7p.m</b>	0.999725	0.999756	0.999756	0.999985	0.999985	0.999985	0.850539	0.429518
<b>8p.m</b>	0.999725	0.999756	0.999756	0.999985	0.999985	0.999985	0.850682	0.429518
<b>9p.m</b>	0.999725	0.999756	0.999756	0.999985	0.999985	0.999985	0.850682	0.429633
<b>10p.m</b>	0.999725	0.999756	0.999756	0.999985	0.999985	0.999985	0.859022	0.429797
<b>11p.m</b>	0.997925	0.999756	0.999985	0.999985	0.999985	0.999985	0.859022	0.429797

Table 3: Daily Bandwidth Utilization for 150 Nodes

<b>Period</b>	<b>@ 64Kbps</b>	<b>@128Kbps</b>	<b>@256Kbps</b>	<b>@512Kbps</b>	<b>@1Mbps</b>	<b>@2Mbps</b>	<b>@4Mbps</b>	<b>@8Mbps</b>
<b>12 midnight</b>	0.941162	0.977325	0.993805	0.997353	0.998554	0.99967	0.999985	0.592067
<b>1a.m</b>	0.941162	0.977325	0.99678	0.999527	0.998554	0.99967	0.999985	0.600257
<b>2a.m</b>	0.967712	0.977325	0.99678	0.999527	0.998554	0.99967	0.999985	0.608869
<b>3a.m</b>	0.967712	0.990601	0.999756	0.99987	0.998554	0.99967	0.999985	0.619426
<b>4a.m</b>	0.967712	0.99472	0.999756	0.99987	0.998554	0.999928	0.999985	0.621536
<b>5a.m</b>	0.967712	0.99472	0.999756	0.99987	0.998554	0.999928	0.999985	0.626686
<b>6a.m</b>	0.967712	0.99472	0.999756	0.99987	0.998554	0.999928	0.999985	0.631907
<b>7a.m</b>	0.992432	0.99884	0.999756	0.99987	0.998955	0.999928	0.999985	0.632436
<b>8a.m</b>	0.992432	0.99884	0.999756	0.99987	0.998955	0.999928	0.999999	0.634131
<b>9a.m</b>	0.992432	0.99884	0.999756	0.99987	0.998955	0.999928	0.999999	0.637622
<b>10a.m</b>	0.992432	0.999756	0.999756	0.99987	0.999298	0.999928	0.999999	0.638859
<b>11a.m</b>	0.992432	0.999756	0.999756	0.99987	0.999298	0.999928	0.999999	0.65006
<b>12noon</b>	0.992432	0.999756	0.999756	0.99987	0.999641	0.999956	0.999999	0.652299
<b>1p.m</b>	0.992432	0.999756	0.999756	0.99985	0.999641	0.999956	0.999999	0.654151
<b>2p.m</b>	0.992432	0.999756	0.999985	0.999985	0.999641	0.999956	0.999999	0.658264
<b>3p.m</b>	0.992432	0.999756	0.999985	0.999985	0.999641	0.999956	0.999999	0.659716
<b>4p.m</b>	0.99884	0.999756	0.999985	0.999985	0.999641	0.999956	0.999999	0.659716
<b>5p.m</b>	0.99884	0.999756	0.999985	0.999985	0.999641	0.999956	0.999999	0.660138
<b>6p.m</b>	0.99884	0.999756	0.999985	0.999985	0.999641	0.999956	0.999999	0.66021
<b>7p.m</b>	0.99884	0.999756	0.999985	0.999985	0.999641	0.999956	0.999999	0.660388
<b>8p.m</b>	0.99884	0.999756	0.999985	0.999985	0.999641	0.999956	0.999999	0.662141
<b>9p.m</b>	0.99884	0.999756	0.999985	0.999985	0.999641	0.999956	0.999999	0.662277
<b>10p.m</b>	0.99884	0.999756	0.999985	0.999985	0.999641	0.999956	0.999999	0.662649
<b>11p.m</b>	0.99884	0.999756	0.999985	0.999985	0.999641	0.999956	0.999999	0.663006

**Table 4: Daily Bandwidth Utilization for 200 Nodes**

<b>Period</b>	<b>@ 64Kbps</b>	<b>@128Kbps</b>	<b>@256Kbps</b>	<b>@512Kbps</b>	<b>@1Mbps</b>	<b>@2Mbps</b>	<b>@4Mbps</b>	<b>@8Mbps</b>
<b>12 midnight</b>	0.90271	0.996094	0.997467	0.992088	0.997009	0.999784	0.999999	0.792196
<b>1a.m</b>	0.983276	0.996094	0.997467	0.995293	0.999126	0.999784	0.999999	0.828001
<b>2a.m</b>	0.983276	0.996094	0.997467	0.999298	0.999928	0.999813	0.999999	0.832279
<b>3a.m</b>	0.983276	0.996094	0.997467	0.999298	0.999928	0.999813	0.999999	0.83308
<b>4a.m</b>	0.983276	0.996094	0.997467	0.999298	0.999928	0.999813	0.999999	0.839782
<b>5a.m</b>	0.983276	0.996094	0.997467	0.999298	0.999928	0.999813	0.999999	0.841348
<b>6a.m</b>	0.983276	0.996094	0.997467	0.999298	0.999928	0.999813	0.999999	0.846705
<b>7a.m</b>	0.983276	0.996094	0.999298	0.999298	0.999928	0.999813	0.999999	0.846705
<b>8a.m</b>	0.983276	0.996094	0.999527	0.999298	0.999928	0.999813	0.999999	0.850925
<b>9a.m</b>	0.983276	0.996094	0.999527	0.999298	0.999928	0.999813	0.999999	0.851169
<b>10a.m</b>	0.997925	0.996094	0.999527	0.999298	0.999928	0.999813	0.999999	0.851769
<b>11a.m</b>	0.997925	0.996094	0.999527	0.999298	0.999928	0.999813	0.999999	0.851898
<b>12noon</b>	0.997925	0.998383	0.999527	0.999298	0.999928	0.999813	0.999999	0.85335
<b>1p.m</b>	0.997925	0.998383	0.999527	0.999298	0.999928	0.999928	0.999999	0.853379
<b>2p.m</b>	0.997925	0.999298	0.999985	0.999298	0.999928	0.999956	0.999999	0.853894
<b>3p.m</b>	0.997925	0.999298	0.999985	0.999298	0.999928	0.999956	0.999999	0.854974
<b>4p.m</b>	0.997925	0.999298	0.999985	0.999413	0.999928	0.999956	0.999999	0.859838
<b>5p.m</b>	0.997925	0.999298	0.999985	0.999413	0.999928	0.999956	0.999999	0.870073
<b>6p.m</b>	0.997925	0.999298	0.999985	0.999413	0.999928	0.999956	0.999999	0.872583
<b>7p.m</b>	0.997925	0.999298	0.999985	0.99987	0.999928	0.999956	0.999999	0.872626
<b>8p.m</b>	0.997925	0.999298	0.999985	0.99987	0.999928	0.999956	0.999999	0.875316
<b>9p.m</b>	0.997925	0.999298	0.999985	0.99987	0.999928	0.999956	0.999999	0.875316
<b>10p.m</b>	0.997925	0.999298	0.999985	0.99987	0.999928	0.999956	0.999999	0.875316
<b>11p.m</b>	0.997925	0.999298	0.999985	0.99987	0.999928	0.999956	0.999999	0.875316

**Table 5: Daily Bandwidth Utilization for 250 Nodes**

<b>Period</b>	<b>@ 64Kbps</b>	<b>@128Kbps</b>	<b>@256Kbps</b>	<b>@512Kbps</b>	<b>@1Mbps</b>	<b>@2Mbps</b>	<b>@4Mbps</b>	<b>@8Mbps</b>
<b>12 midnight</b>	<b>0.996094</b>	<b>0.991974</b>	<b>0.986023</b>	<b>0.997696</b>	<b>0.999355</b>	<b>0.999269</b>	<b>0.999556</b>	<b>0.999899</b>
<b>1a.m</b>	<b>0.996094</b>	<b>0.991974</b>	<b>0.996323</b>	<b>0.999985</b>	<b>0.999355</b>	<b>0.99928</b>	<b>0.999641</b>	<b>0.999899</b>
<b>2a.m</b>	<b>0.996094</b>	<b>0.991974</b>	<b>0.999298</b>	<b>0.999985</b>	<b>0.999355</b>	<b>0.99928</b>	<b>0.999756</b>	<b>0.999963</b>
<b>3a.m</b>	<b>0.996094</b>	<b>0.991974</b>	<b>0.999298</b>	<b>0.999985</b>	<b>0.999355</b>	<b>0.99928</b>	<b>0.999928</b>	<b>0.999978</b>
<b>4a.m</b>	<b>0.996094</b>	<b>0.995178</b>	<b>0.999298</b>	<b>0.999985</b>	<b>0.999355</b>	<b>0.99928</b>	<b>0.999985</b>	<b>0.999978</b>
<b>5a.m</b>	<b>0.996094</b>	<b>0.995178</b>	<b>0.999298</b>	<b>0.999985</b>	<b>0.999355</b>	<b>0.99928</b>	<b>0.999985</b>	<b>0.999978</b>
<b>6a.m</b>	<b>0.996094</b>	<b>0.995178</b>	<b>0.999298</b>	<b>0.999985</b>	<b>0.999355</b>	<b>0.99928</b>	<b>0.999985</b>	<b>0.999999</b>
<b>7a.m</b>	<b>0.996094</b>	<b>0.998383</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.999355</b>	<b>0.99928</b>	<b>0.999985</b>	<b>0.999999</b>
<b>8a.m</b>	<b>0.996094</b>	<b>0.998383</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.999641</b>	<b>0.999928</b>	<b>0.999985</b>	<b>0.999999</b>
<b>9a.m</b>	<b>0.996094</b>	<b>0.998383</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.999641</b>	<b>0.999928</b>	<b>0.999985</b>	<b>0.999999</b>
<b>10a.m</b>	<b>0.996094</b>	<b>0.998383</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.999641</b>	<b>0.999928</b>	<b>0.999985</b>	<b>0.999999</b>
<b>11a.m</b>	<b>0.996094</b>	<b>0.998383</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.999641</b>	<b>0.999928</b>	<b>0.999985</b>	<b>0.999999</b>
<b>12noon</b>	<b>0.996094</b>	<b>0.998383</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.999641</b>	<b>0.999928</b>	<b>0.999985</b>	<b>0.999999</b>
<b>1p.m</b>	<b>0.996094</b>	<b>0.998383</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.999756</b>	<b>0.999956</b>	<b>0.999985</b>	<b>0.999999</b>
<b>2p.m</b>	<b>0.996094</b>	<b>0.998383</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.999756</b>	<b>0.999956</b>	<b>0.999985</b>	<b>0.999999</b>
<b>3p.m</b>	<b>0.996094</b>	<b>0.998383</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.999756</b>	<b>0.999956</b>	<b>0.999985</b>	<b>0.999999</b>
<b>4p.m</b>	<b>0.996094</b>	<b>0.998383</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.999756</b>	<b>0.999956</b>	<b>0.999985</b>	<b>0.999999</b>
<b>5p.m</b>	<b>0.996094</b>	<b>0.99884</b>	<b>0.999756</b>	<b>0.999985</b>	<b>0.999756</b>	<b>0.999956</b>	<b>0.999985</b>	<b>0.999999</b>
<b>6p.m</b>	<b>0.996094</b>	<b>0.99884</b>	<b>0.999985</b>	<b>0.999985</b>	<b>0.999756</b>	<b>0.999956</b>	<b>0.999985</b>	<b>0.999999</b>
<b>7p.m</b>	<b>0.996094</b>	<b>0.99884</b>	<b>0.999985</b>	<b>0.999985</b>	<b>0.999756</b>	<b>0.999956</b>	<b>0.999985</b>	<b>0.999999</b>
<b>8p.m</b>	<b>0.996094</b>	<b>0.99884</b>	<b>0.999985</b>	<b>0.999985</b>	<b>0.999756</b>	<b>0.999956</b>	<b>0.999985</b>	<b>0.999999</b>
<b>9p.m</b>	<b>0.996094</b>	<b>0.99884</b>	<b>0.999985</b>	<b>0.999985</b>	<b>0.999756</b>	<b>0.999956</b>	<b>0.999985</b>	<b>0.999999</b>
<b>10p.m</b>	<b>0.996094</b>	<b>0.99884</b>	<b>0.999985</b>	<b>0.999985</b>	<b>0.999756</b>	<b>0.999956</b>	<b>0.999985</b>	<b>0.999999</b>
<b>11p.m</b>	<b>0.996094</b>	<b>0.99884</b>	<b>0.999985</b>	<b>0.999985</b>	<b>0.999756</b>	<b>0.999956</b>	<b>0.999999</b>	<b>0.999999</b>

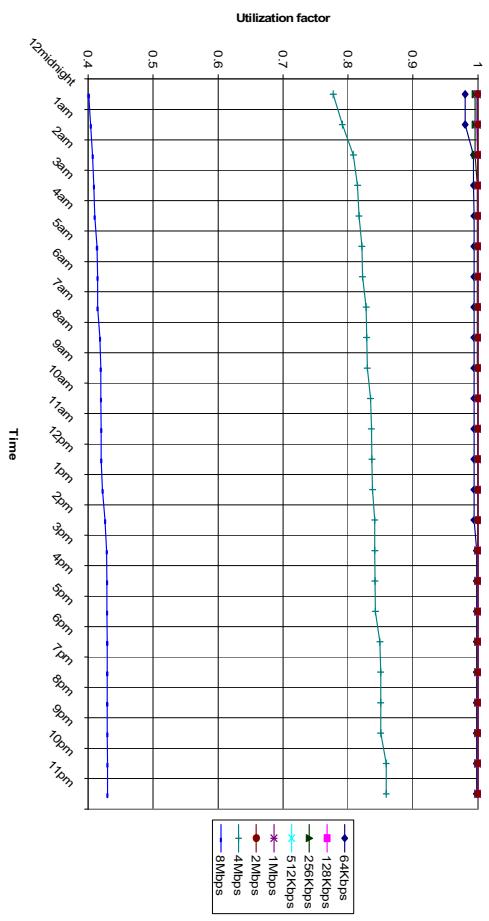


FIGURE 2: DAILY BANDWIDTH UTILIZATION FOR 10 NODES

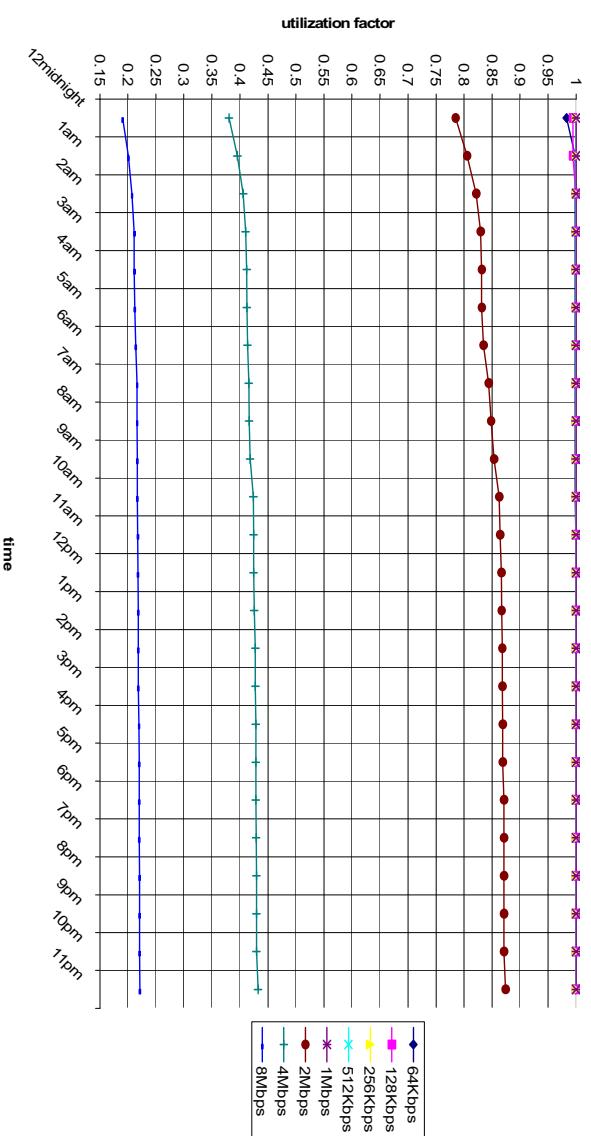


FIGURE 1: DAILY BANDWIDTH UTILIZATION FOR 50 NODES

FIGURE 3: DAILY BANDWIDTH UTILIZATION FOR 150 NODES 0 NODES

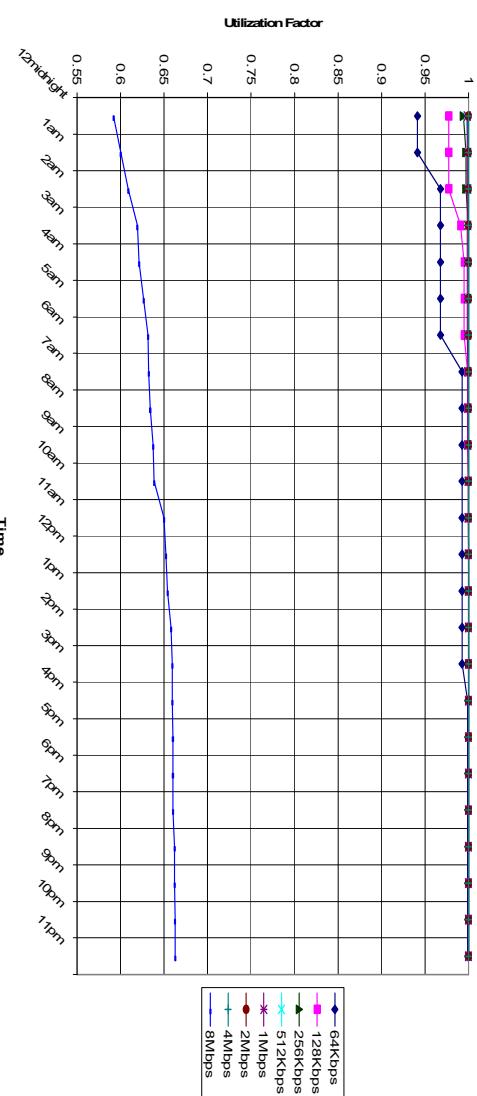
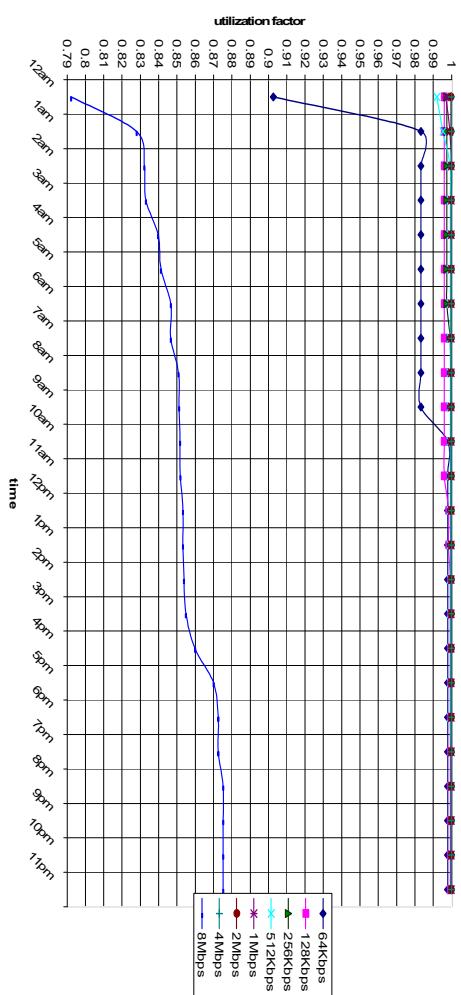


FIGURE 4: DAILY BANDWIDTH UTILIZATION FOR 200 NODES



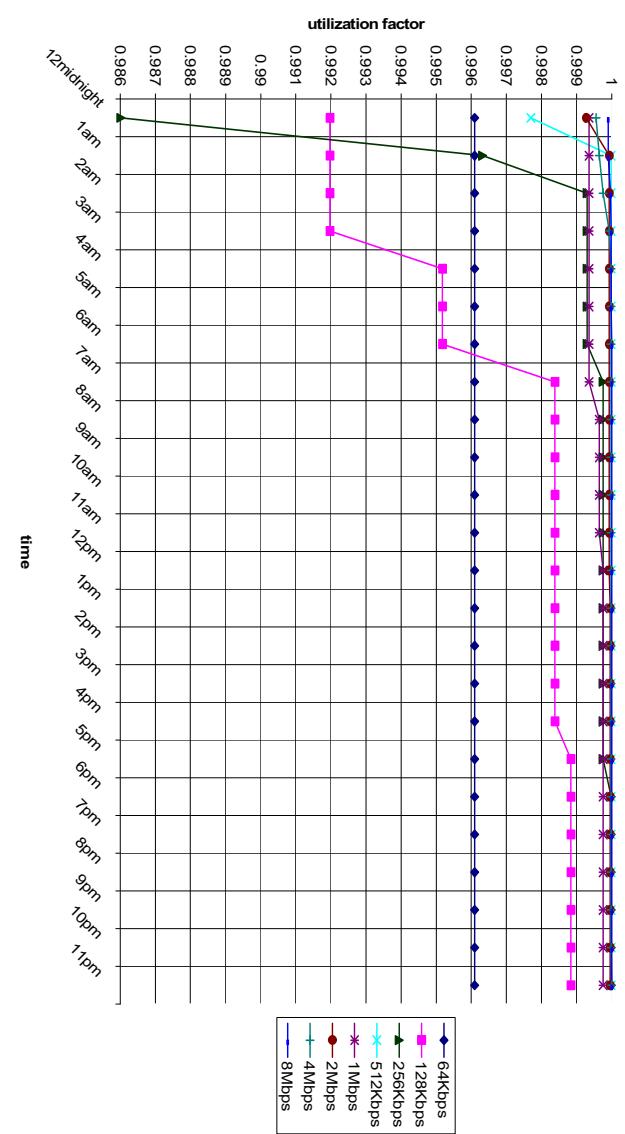


FIGURE 5: DAILY BANDWIDTH UTILIZATION FOR 250 NODES

## INFERENCES FROM TABLES AND FIGURES

From the tables and the graphs (tables 1 - 5 and figures 1 – 5), the following inferences could be inferred for various nodes at different bandwidths.

- (i) The bandwidth consumption varies randomly with time and the number of nodes. The arrival pattern portrays the peak periods to be varying among the nodes.
- (ii) With 50 nodes the daily bandwidth consumption curve is fairly or approximately linear. With 150, 200 and 250 nodes it gets to a point where there is a considerable increase in the bandwidth consumption and the curve becomes non-linear. The implication is that bandwidth is wasted whenever there is no proportional increment in available nodes.
- (iii) There is always a critical point in bandwidth availability at which further increase in available bandwidth (BW) did not improve the utilization factor for the nodes. In actual fact the utilization factor starts decreasing. If the cost of purchase of bandwidth is to be minimized and available bandwidth optimized, VPN must not be operated above the point.
- (iv) Each node acts adaptively and optimally to the dynamics of the external environment so that the available bandwidths are shared optimally for each node despite the fact that each node behaves as a selfish node.
- (v) The utilization factor can be as high as 99.99% (250 nodes @ 4096/8192kbps).

## CONCLUSION

Simulation results show that if the cost of purchase of bandwidth is to be minimized and the available bandwidth optimized, the number of nodes and utility must be commensurate with the quantity of bandwidth purchased by operators of Virtual Private Network. The utilization factors were as high as ninety nine point ninety nine percent (99.99%).

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